

Disjoint Paths in Graphs III, Characterization

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Abstract

Let G be a graph, $\{a, b, c\} \subseteq V(G)$, and $\{a', b', c'\} \subseteq V(G)$ such that $\{a, b, c\} \neq \{a', b', c'\}$. We say that $(G, \{a, b, c\}, \{a', b', c'\})$ is an *obstruction* if, for any three vertex disjoint paths from $\{a, b, c\}$ to $\{a', b', c'\}$ in G , one path is from b to b' . In this paper we characterize obstructions.

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1 Introduction

We use the terminology in [5]. Let G be a graph, $\{a, b, c\} \subseteq V(G)$, and $\{a', b', c'\} \subseteq V(G)$. Then $(G, \{a, b, c\}, \{a', b', c'\}, (b, b'))$ is an *obstruction* if for any three vertex disjoint paths from $\{a, b, c\}$ to $\{a', b', c'\}$ in G , one path is from b to b' . In [6], a special class of obstructions are characterized. In this paper, we characterize all obstructions. In order to state our main result, we need the following definition from [5].

A *3-planar graph* (G, \mathcal{A}) consists of a graph G and a set $\mathcal{A} = \{A_1, \dots, A_k\}$ of pairwise disjoint subsets of $V(G)$ (possibly $\mathcal{A} = \emptyset$) such that

- (a) for $1 \leq i \neq j \leq k$, $N(A_i) \cap A_j = \emptyset$,
- (b) for $1 \leq i \leq k$, $|N(A_i)| \leq 3$, and
- (c) if $p(G, \mathcal{A})$ denotes the graph obtained from G by (for each i) deleting A_i and adding edges joining every pair of distinct vertices in $N(A_i)$, then $p(G, \mathcal{A})$ can be drawn in a closed disc D with no pair of edges crossing such that, for each A_i with $|N(A_i)| = 3$, $N(A_i)$ induces a facial triangle in $p(G, \mathcal{A})$.

If, in addition, b_0, b_1, \dots, b_n are vertices of G such that $b_i \notin A_j$ for any $A_j \in \mathcal{A}$ and b_0, b_1, \dots, b_n occur on the boundary of D in that cyclic order, then we say that $(G, \mathcal{A}, b_0, b_1, \dots, b_n)$ is *3-planar*. If there is no need to specify \mathcal{A} , we will simply say that $(G, b_0, b_1, \dots, b_n)$ is 3-planar.

The building blocks of obstructions are described in the following definition.

(1.1) Definition. Let G be a graph, $\{a, b, c\} \subseteq V(G)$, and $\{a', b', c'\} \subseteq V(G)$. Suppose $\{a, b, c\} \neq \{a', b', c'\}$, and assume that G has no 3-separation (G_1, G_2) such that $\{a, b, c\} \subseteq G_1$ and $\{a', b', c'\} \subseteq G_2$. Then we call $(G, (a, b, c), (a', b', c'))$ a *rung* if one of the following conditions is satisfied:

- (1) $b = b'$ or $\{a, c\} = \{a', c'\}$;
- (2) $a = a'$ and $(G - a, c, c', b', b)$ is 3-planar;
- (2') $c = c'$ and $(G - c, a, a', b', b)$ is 3-planar;
- (3) $\{a, b, c\} \cap \{a', b', c'\} = \emptyset$ and (G, a', b', c', c, b, a) is 3-planar;
- (4) $\{a, b, c\} \cap \{a', b', c'\} = \emptyset$, G has a 1-separation (G_1, G_2) such that $\{a, a', b, b'\} \subseteq G_1$, $\{c, c'\} \subseteq G_2$, and (G_1, a, a', b', b) is 3-planar;
- (4') $\{a, b, c\} \cap \{a', b', c'\} = \emptyset$, G has a 1-separation (G_1, G_2) such that $\{c, c', b, b'\} \subseteq G_1$, $\{a, a'\} \subseteq G_2$, and (G_1, c, c', b', b) is 3-planar;
- (5) $\{a, b, c\} \cap \{a', b', c'\} = \emptyset$, (G, a, a', b', b) is 3-planar, and G has a separation (G_1, G_2) such that $V(G_1 \cap G_2) = \{z, b\}$ (or $V(G_1 \cap G_2) = \{z, b'\}$), $\{a, a', b, b'\} \subseteq G_1$, $\{c, c'\} \subseteq G_2$, and (G_2, c, c', z, b) (or (G_2, c, c', b', z)) is 3-planar;

- (5') $\{a, b, c\} \cap \{a', b', c'\} = \emptyset$, (G, c, c', b', b) is 3-planar, and G has a separation (G_1, G_2) such that $V(G_1 \cap G_2) = \{z, b\}$ (or $V(G_1 \cap G_2) = \{z, b'\}$), $\{c, c', b, b'\} \subseteq G_1$, $\{a, a'\} \subseteq G_2$, and (G_2, a, a', z, b) (or (G_2, a, a', b', z)) is 3-planar;
- (6) $\{a, b, c\} \cap \{a', b', c'\} = \emptyset$, and there are pairwise edge disjoint subgraphs G_a, G_c, M of G such that $G = G_a \cup G_c \cup M$, $V(G_a \cap M) = \{u, w\}$, $V(G_c \cap M) = \{p, q\}$, $V(G_a \cap G_c) = \emptyset$, $\{a, a', b'\} \subseteq G_a$, $\{c, c', b\} \subseteq G_c$, (G_a, a, a', b', w, u) is 3-planar, and (G_c, c', c, b, p, q) is 3-planar;
- (6') $\{a, b, c\} \cap \{a', b', c'\} = \emptyset$, and there are pairwise edge disjoint subgraphs G_a, G_c, M of G such that $G = G_a \cup G_c \cup M$, $V(G_a \cap M) = \{u, w\}$, $V(G_c \cap M) = \{p, q\}$, $V(G_a \cap G_c) = \emptyset$, $\{a, a', b\} \subseteq G_a$, $\{c, c', b'\} \subseteq G_c$, (G_a, b, a, a', w, u) is 3-planar, and (G_c, b', c', c, p, q) is 3-planar;
- (7) $\{a, b, c\} \cap \{a', b', c'\} = \emptyset$, and there are pairwise edge disjoint subgraphs G_a, G_c, M of G such that $G = G_a \cup G_c \cup M$, $V(G_a \cap M) = \{b, b', w\}$, $V(G_c \cap M) = \{b, b', p\}$, $V(G_a \cap G_c) = \{b, b'\}$, $\{a, a'\} \subseteq G_a$, $\{c, c'\} \subseteq G_c$, (G_a, a, a', b', w, b) is 3-planar, and (G_c, c', c, b, p, b') is 3-planar.

(1.2) Definition. Let L be a graph and let R_1, \dots, R_m be edge disjoint subgraphs of L such that

- (i) $(R_i, (x_{i-1}, v_{i-1}, y_{i-1}), (x_i, v_i, y_i))$ is a rung for each $1 \leq i \leq m$,
- (ii) $V(R_i \cap R_j) = \{x_i, v_i, y_i\} \cap \{x_{j-1}, v_{j-1}, y_{j-1}\}$ for $1 \leq i < j \leq m$,
- (iii) for any $1 \leq i < j \leq m$, if $x_i = x_j$ then $x_k = x_i$ for all $i \leq k \leq j$, if $v_i = v_j$ then $v_k = v_i$ for all $i \leq k \leq j$, and if $y_i = y_j$ then $y_k = y_i$ for all $i \leq k \leq j$.
- (iv) $L = (\bigcup_{i=1}^m R_i) + S$, where S consists of edges of L with both ends in some $\{x_i, v_i, y_i\}$, $1 \leq i \leq m$.

Then we call $(L, (x_0, v_0, y_0), (x_m, v_m, y_m))$ a *ladder with rungs* $(R_i, (x_{i-1}, v_{i-1}, y_{i-1}), (x_i, v_i, y_i))$, $i = 1, \dots, m$, or simply, a *ladder along* $v_0 \dots v_m$.

It is easy to see that a rung is at most 5-connected. It was shown in [5] (Proposition 4.4) that if $(L, (x_0, v_0, y_0), (x_m, v_m, y_m))$ is a ladder, then $(L, \{x_0, y_0\}, \{x_m, y_m\}, (v_0, v_m))$ is an obstruction. The main result of this paper states that every obstruction can be constructed from ladders and 3-planar graphs in a special way. For a sequence S , the *reduced sequence* of S is the sequence obtained from S by removing all but one consecutive identical elements. For example, the reduced sequence of $aaabcca$ is $abca$.

(1.3) Theorem. Let G be a graph, $\{a, b, c\} \subseteq V(G)$, and $\{a', b', c'\} \subseteq V(G)$ such that $\{a, b, c\} \neq \{a', b', c'\}$. Assume that, for any $T \subseteq V(G)$ with $|T| \leq 3$, every component of $G - T$ contains some element of $\{a, b, c\} \cup \{a', b', c'\}$. Then $(G, \{a, c\}, \{a', c'\}, (b, b'))$ is an obstruction iff one of the following statements holds.

- (1) G has a separation (G_1, G_2) of order at most 2 such that $\{a, b, c\} \subseteq V(G_1)$ and $\{a', b', c'\} \subseteq V(G_2)$.
- (2) $(G, (a, b, c), (a', b', c'))$ is a ladder.
- (3) G has a separation (J, L) such that $V(J \cap L) = \{w_0, \dots, w_n\}$, (J, w_0, \dots, w_n) is 3-planar, $(L, (a, b, c), (a', b', c'))$ is a ladder along a sequence $v_0 \dots v_m$, where $v_0 = b$, $v_m = b'$, and $w_0 \dots w_n$ is the reduced sequence of $v_0 \dots v_m$.

Note the condition of (1.3) that for any $T \subseteq V(G)$ with $|T| \leq 3$, every component of $G - T$ contains some element of $\{a, b, c\} \cup \{a', b', c'\}$. It is a natural condition, for the following reason. Suppose that $T \subseteq V(G)$, $|T| \leq 3$, and $G - T$ contains a component H with $V(H) \cap (\{a, b, c\} \cup \{a', b', c'\}) = \emptyset$. Let G' be obtained from G by removing H and adding new edges between every pair of distinct vertices in $N(H)$. Then it is easy to see that $(G, \{a, c\}, \{a', c'\}, (b, b'))$ is an obstruction iff $(G', \{a, c\}, \{a', c'\}, (b, b'))$ is an obstruction.

Also note that if (1) of (1.3) holds, then there do not exist three disjoint paths from $\{a, b, c\}$ to $\{a', b', c'\}$, and hence, $(G, \{a, c\}, \{a', c'\}, (b, b'))$ is an obstruction in a trivial sense.

As a consequence of (1.3), we will prove the following result.

(1.4) Corollary. *Let $(G, \{a, c\}, \{a', c'\}, (b, b'))$ be an obstruction. If $\{a, c\} \neq \{a', c'\}$ and $b \neq b'$, then G is at most 7-connected.*

The rest of the paper is organized as follows. In Section 2, we prove a technical lemma. In Section 3, we prove (1.3). In Section 4, we will prove (1.4) and construct 7-connected obstructions.

2 Good ladders

The main goal of this section is to prove a technical lemma, based on good ladders. To this end, we need the following result proved in [2] (also independently in [1] and [3]). Also, see (2.4) of [5].

(2.1) Theorem. *Let G be a graph and $\{s_1, s_2, t_1, t_2\} \subseteq V(G)$. Then G contains no $(\{s_1, t_1\}, \{s_2, t_2\})$ -linkage iff (G, s_1, s_2, t_1, t_2) is 3-planar.*

We also need the following result, which follows easily from Proposition 3.2 of [5].

(2.2) Lemma. *Let (G, \mathcal{A}) be 3-planar and let $b, b' \in V(G)$ be distinct. Let $u, v \in V(G)$ such that $u, v \in p(G, \mathcal{A})$. If $p(G, \mathcal{A})$ contains a $(\{u, b\}, \{v, b'\})$ -linkage, then G contains a $(\{u, b\}, \{v, b'\})$ -linkage L .*

To prove the main result of this section, we need the main result of [6]. For convenience, we state the following definition.

(2.3) Definition. Let $(G, (a, b, c), (a', b', c'))$ be a rung. Suppose that $G - \{b, b'\}$ has disjoint paths A, C from a, c to a', c' , respectively, and assume that the following conditions are satisfied:

- (1) If G is connected and $b \neq b'$, then $G - ((A - a') \cup (C - c))$ contains a $(\{b, a'\}, \{c, b'\})$ -linkage, and $G - ((A - a) \cup (C - c'))$ contains an $(\{a, b'\}, \{b, c'\})$ -linkage;
- (1') if G is connected and $b = b'$, then $G - ((A - a) \cup C)$ has a path from a to b' , $G - ((A - a') \cup C)$ has a path from a' to b , $G - (A \cup (C - c))$ has a path from c to b' , and $G - (A \cup (C - c'))$ has a path from c' to b ;
- (2) if G is not connected and $\{a, b\} \cup \{a', b'\}$ is contained in a component of G , then $G - ((A - a) \cup C)$ contains a path from a to b' (and not using b if $b \neq b'$) and $G - ((A - a') \cup C)$ contains a path from b to a' (and not using b' if $b \neq b'$); and
- (2') if G is not connected and $\{b, c\} \cup \{b', c'\}$ is contained in a component of G , then $G - (A \cup (C - c))$ contains a path from c to b' (and not using b if $b \neq b'$), and $G - (A \cup (C - c'))$ contains a path from b to c' (and not using b' if $b \neq b'$).

Then we call $(G, (a, b, c), (a', b', c'), A, C)$ a *good rung*. Let $(L, (x_0, v_0, y_0), (x_m, v_m, y_m))$ be a ladder with good rungs $(R_i, (x_{i-1}, v_{i-1}, y_{i-1}), (x_i, v_i, y_i), X_i, Y_i)$, $i = 1, \dots, m$. Then we say that $(L, (x_0, v_0, y_0), (x_m, v_m, y_m), X, Y)$ is a *good ladder* along $v_0 \dots v_m$, where $X = \bigcup_{i=1}^m X_i$ and $Y = \bigcup_{i=1}^m Y_i$.

The following is the main result of [6].

(2.4) Lemma. Let $(G, \{a, c\}, \{a', c'\}, (b, b'))$ be an obstruction satisfying the following conditions:

- (i) $G - \{b, b'\}$ contains disjoint paths A, C from a, c to a', c' , respectively,
- (ii) $A \cup C$ is an induced subgraph of G , and
- (iii) $G - (A \cup C)$ is connected.

Then one of the following statements holds:

- (1) There is a separation (G_1, G_2) in G of order at most 2 such that $\{a, b, c\} \subseteq V(G_1)$ and $\{a', b', c'\} \subseteq V(G_2)$.
- (2) There is a subset $T \subseteq V(G)$ such that $|T| \leq 3$ and some component of $G - T$ contains no element of $\{a, b, c\} \cup \{a', b', c'\}$.
- (3) $(G, (a, b, c), (a', b', c'), A, C)$ is a good ladder.
- (4) G has a separation (J, L) such that $V(J \cap L) = \{w_0, \dots, w_n\}$, (J, w_0, \dots, w_n) is 3-planar, $(L, (a, b, c), (a', b', c'), A, C)$ is a good ladder along a sequence $v_0 \dots v_m$, where $v_0 = b$, $v_m = b'$, and $w_0 \dots w_n$ is the reduced sequence of $v_0 \dots v_m$.

We can now state and prove the technical lemma mentioned earlier. The conditions of the lemma arise in the proof of (1.3).

(2.5) Lemma. *Let $(L, (a, b, c), (a', b', c'), A, C)$ be a good ladder with good rungs $(R_i, (x_{i-1}, v_{i-1}, y_{i-1}), (x_i, v_i, y_i), A_i, C_i)$, where $A_i = A \cap R_i$, $C_i = C \cap R_i$, $a = x_0$, $a' = x_m$, $b = v_0$, $b' = v_m$, $c = y_0$, and $c' = y_m$. Let L' be a graph and let L^* be obtained from the disjoint union of L and L' by adding edges between $V(A \cup C)$ and $V(L')$ and between $V(A)$ and $V(C)$. Assume the following conditions hold. (We use N instead of N_{L^*} to denote neighborhood in L^* .)*

- (a) *For any $i, j \in \{1, \dots, m\}$ with $i \leq j$, let $L_{i,j}$ denote the graph obtained from L^* by deleting all $R_k - (A_k \cup C_k \cup \{v_{i-1}, v_j\})$ with $k < i$ or $k > j$ and by deleting v_0, \dots, v_m except v_{i-1} and v_j . Then $(L_{i,j}, \{a, c\}, \{a', c'\}, (v_{i-1}, v_j))$ is an obstruction.*
- (b) *If $q \in V(A)$ (respectively, $q \in V(C)$) such that L contains a path from q to $\{v_0, \dots, v_m\}$ internally disjoint from $V(A \cup C) \cup \{v_0, \dots, v_m\}$, then, for any component D of L' , $N(D) \cap V(A) \subseteq aAq$ or $N(D) \cap V(A) \subseteq qAa'$ (respectively, $N(D) \cap V(C) \subseteq cCq$ or $N(D) \cap V(C) \subseteq qCc'$).*
- (c) *If L^* contains paths from distinct $s, s' \in V(A)$ (in that order from a to a') to distinct $t, t' \in V(C)$ (in that order from c' to c), respectively, which are internally disjoint from L , then L contains no path from $V(sAs' - \{s, s'\}) \cup V(tCt' - \{t, t'\})$ to $\{v_0, \dots, v_m\}$ internally disjoint from $V(A \cup C) \cup \{v_0, \dots, v_m\}$.*
- (d) *Suppose that L^* contains paths from $s, s' \in V(A)$ (in that order from a to a') to $t, t' \in V(C)$ (in that order from c to c'), respectively, which are internally disjoint from L . If L contains a path from $V(tCt' - \{t, t'\})$ (respectively, $V(sAs' - \{s, s'\})$) to $\{v_0, \dots, v_m\}$ internally disjoint from $V(A \cup C) \cup \{v_0, \dots, v_m\}$, then there is a vertex $q \in sAs'$ (respectively, $q \in tCt'$) such that, for every component D of L' , $N(D) \cap V(A) \subseteq aAq$ or $N(D) \cap V(A) \subseteq qAa'$ (respectively, $N(D) \cap V(C) \subseteq cCq$ or $N(D) \cap V(C) \subseteq qCc'$).*
- (e) *If R_i is not connected, but $\{x_{i-1}, v_{i-1}\} \cup \{x_i, v_i\}$ (respectively, $\{y_{i-1}, v_{i-1}\} \cup \{y_i, v_i\}$) is contained in a component of R_i , then $v_{i-1} = v_i$.*
- (f) *For any $T \subseteq V(L^*)$ with $|T| \leq 3$, every component of $L^* - T$ contains a vertex in $V(A \cup C) \cup \{v_0, \dots, v_m\}$.*

Then $(L^*, (a, b, c), (a', b', c'))$ is a ladder along a sequence $z_0 \dots z_p$, and the reduced sequence of $z_0 \dots z_p$ is the reduced sequence of $v_0 \dots v_m$.

Proof. Suppose (2.5) fails. Choose L^*, L, A, C so that $|V(L^*)| + |E(L^*)|$ is minimum. Then it is easy to see that $\{a, b, c\}$ and $\{a', b', c'\}$ are independent sets in L^* . We proceed by proving Claims 1–6.

Claim 1. L^* does not contain any separation (G_1, G_2) such that $V(G_1 \cap G_2) = \{a'', v_k, c''\}$, $\{a, b, c\} \subseteq G_1$, and $\{a', b', c'\} \subseteq G_2$, where $k \in \{0, \dots, m\}$, $a'' \in V(A)$, and $c'' \in V(C)$.

Suppose that such a separation (G_1, G_2) does exist. Note that $\{a'', v_k, c''\} \neq \{a, b, c\}$ and $\{a'', v_k, c''\} \neq \{a', b', c'\}$; for otherwise, since $\{a, b, c\}$ and $\{a', b', c'\}$ are independent in L^* , $L^* - \{a'', v_k, c''\}$ has a component not containing any element of $V(A) \cup V(C) \cup \{v_0, \dots, v_m\}$, contradicting (f).

For $i \in \{1, 2\}$, let $L_i = L \cap G_i$, $L'_i = L' \cap G_i$, $X_i = A \cap G_i$ and $Y_i = C \cap G_i$.

First, we claim that $(L_1, (a, b, c), (a'', v_k, c''), X_1, Y_1)$ is a good ladder along a sequence whose reduced sequence is the reduced sequence of $v_0 \dots v_k$. This is clear if, for every $i \in \{1, \dots, m\}$, either $R_i \subseteq L_1$ or $R_i \subseteq L_2$. So assume that p is minimum such that $R_p \not\subseteq L_1$ and $R_p \not\subseteq L_2$, and q is maximum such that $R_q \not\subseteq L_1$ and $R_q \not\subseteq L_2$. Then, for each $p \leq i \leq q$, $\{a_{i-1}, v_{i-1}, c_{i-1}\} \not\subseteq L_2$ and $\{a_i, v_i, c_i\} \not\subseteq L_1$, and hence, $R_i \not\subseteq L_1$ and $R_i \not\subseteq L_2$. Let $R_i^1 = R_i$ for $1 \leq i < p$; then $(R_i^1, (x_{i-1}, v_{i-1}, y_{i-1}), (x_i, v_i, y_i), A_i, C_i)$ is a good rung. Now assume $i \in \{p, \dots, q\}$. Because R_i has no 3-separation (R'_i, R''_i) such that $\{x_{i-1}, v_{i-1}, y_{i-1}\} \subseteq R'_i$ and $\{x_i, v_i, y_i\} \subseteq R''_i$, either $x_{i-1}Ax_i \subseteq aAa''$ and $y_{i-1}Cy_i \subseteq c''Cc'$ or $x_{i-1}Ax_i \subseteq a''Aa'$ and $y_{i-1}Cy_i \subseteq cCc''$. This implies that $a'' \in \{x_{p-1}, \dots, x_q\}$ and $c'' \in \{y_{p-1}, \dots, y_q\}$. Let $R_i^1 = (R_i \cap L_1) \cup \{c''\}$ if $x_{i-1}Ax_i \subseteq aAa''$, and let $R_i^1 = (R_i \cap L_1) \cup \{a''\}$ if $y_{i-1}Cy_i \subseteq cCc''$. Note that R_i^1 is not connected. Since $(R_i, (x_{i-1}, v_{i-1}, y_{i-1}), (x_i, v_i, y_i), A_i, C_i)$ is a good rung, it is easy to see that (by using (2.3)) if $x_{i-1}Ax_i \subseteq aAa''$ then $(R_i^1, (x_{i-1}, v_{i-1}, c''), (x_i, v_k, c''), A_i, \{c''\})$ is a good rung, and if $y_{i-1}Cy_i \subseteq cCc''$ then $(R_i^1, (a'', v_{i-1}, y_{i-1}), (a'', v_k, y_i), \{a''\}, C_i)$ is a good rung. Hence, $(L_1, (a, b, c), (a'', v_k, c''), X_1, Y_1)$ is a good ladder along a sequence whose reduced sequence is the reduced sequence of $v_0 \dots v_k$.

Similarly, we can show that $(L_2, (a'', v_k, c''), (a', b', c'), X_2, Y_2)$ is a good ladder along a sequence whose reduced sequence is the reduced sequence of $v_k \dots v_m$.

It is easy to see that, for $i \in \{1, 2\}$, G_i, L_i, L'_i, X_i, Y_i (as L^*, L, L', A, C , respectively) also satisfy (a)–(f) of (2.5) (with a'', v_k, c'' as a', b', c' , respectively, when $i = 1$, and with a'', b'', c'' as a, b, c , respectively, when $i = 2$).

Since $|V(G_1)| + |E(G_1)| < |V(L^*)| + |E(G^*)|$, $(G_1, (a, b, c), (a'', v_k, c''))$ is a ladder along $z_0 \dots z_q$, where $z_0 = v_0$ and $z_q = v_k$, and the reduced sequence of $z_0 \dots z_q$ is the reduced sequence of $v_0 \dots v_k$. Similarly, $(G_2, (a'', v_k, c''), (a', b', c'))$ is a ladder along $z_q \dots z_p$, where $z_q = v_k$ and $z_p = v_m$, and the reduced sequence of $z_q \dots z_p$ is the reduced sequence of $v_k \dots v_m$. Hence, $(L^*, (a, b, c), (a', b', c'))$ is a ladder along $z_0 \dots z_p$, where $z_0 = v_0 = b$ and $z_p = v_m = b'$, and the reduced sequence of $z_0 \dots z_p$ is the reduced sequence of $v_0 \dots v_m$. This is a contradiction.

Claim 2. $b \neq b'$ and $\{a, c\} \neq \{a', c'\}$.

Suppose $b = b'$ or $\{a, c\} = \{a', c'\}$. Then by Claim 1, $(L^*, (a, b, c), (a', b', c'))$ is a rung (as in (1) of (1.1)), and so, a ladder, contradicting the choice of L^* .

Claim 3. (1) If $\{x_{i-1}, v_{i-1}\} \cup \{x_i, v_i\}$ is contained in a component of R_i , then for each $x \in \{x_{i-1}, x_i\}$ and for every component D of L' , $N(D) \cap V(A) \subseteq aAx$ or $N(D) \cap V(A) \subseteq xAa'$. (2) If $\{y_{i-1}, v_{i-1}\} \cup \{y_i, v_i\}$ is contained in a component of R_i , then for each $y \in \{y_{i-1}, y_i\}$ and for every component D of L' , $N(D) \cap V(C) \subseteq cCy$ or $N(D) \cap V(C) \subseteq yCc'$.

Assume that $\{x_{i-1}, v_{i-1}\} \cup \{x_i, v_i\}$ is contained in a component of R_i . Because $(R_i, (x_{i-1}, v_{i-1}, y_{i-1}), (x_i, v_i, y_i), A_i, C_i)$ is a good rung, it follows from (1) or (2) of (2.3) that, for each $x \in \{x_{i-1}, x_i\}$, $R_i - ((A - x) \cup C)$ contains a path from x to $\{v_{i-1}, v_i\}$. Hence, (1) follows from (b) (with x as $q \in V(A)$). Similarly, we can prove (2).

Claim 4. If R_i is connected then $v_{i-1} = v_i$.

Suppose R_i is connected and $v_{i-1} \neq v_i$. Since $(R_i, (x_{i-1}, v_{i-1}, y_{i-1}), (x_i, v_i, y_i), A_i, C_i)$ is a good rung, it follows from (1) of (2.3) that $R_i - ((A_i - x_i) \cup (C_i - y_{i-1}))$ contains a $(\{v_{i-1}, x_i\}, \{y_{i-1}, v_i\})$ -linkage X and $R_i - ((A_i - x_{i-1}) \cup (C_i - y_i))$ contains a $(\{v_{i-1}, y_i\}, \{x_{i-1}, v_i\})$ -linkage Y .

Then L^* has no path from $aAx_i - x_i$ to $y_{i-1}Cc' - y_{i-1}$ internally disjoint from L . Otherwise, let P be a path from $p \in V(aAx_i - x_i)$ to $p' \in V(y_{i-1}Cc' - y_{i-1})$ internally disjoint from L . Let $K = aAp \cup P \cup p'Cc' \cup cCy_{i-1} \cup x_iAa' \cup X$. Then, $K \subseteq L_{i,i}$, and K contains disjoint paths from a, v_{i-1}, c to c', a', v_i , respectively, contradicting (a).

Similarly, we can use Y (instead of X) to show that L^* has no path from $x_{i-1}Aa' - x_{i-1}$ to $cCy_i - y_i$ internally disjoint from L .

Therefore, $\{x_{i-1}, v_{i-1}, y_{i-1}\} = \{a, b, c\}$; for otherwise, by Claim 3, L^* has a separation (G_1, G_2) such that $V(G_1 \cap G_2) = \{x_{i-1}, v_{i-1}, y_{i-1}\}$, $\{a, b, c\} \subseteq G_1$, and $\{a', b', c'\} \subseteq G_2$, contradicting Claim 1. Similarly, we can prove $\{x_i, v_i, y_i\} = \{a', b', c'\}$. Hence by Claim 2, bb' is the reduced sequence of $v_0 \dots v_m$.

If $L' = \emptyset$, then $L^* = R_i$, and so, $(L^*, (a, b, c), (a', b', c'))$ is a ladder along bb' , contradicting the choice of L^* .

So assume that $L' \neq \emptyset$. Let D be a component of L' . By (f), $|N(D) \cap V(A)| \geq 2$ or $|N(D) \cap V(C)| \geq 2$. By symmetry, let $u, v \in V(A) \cap N(D)$ with $u \neq v$. By (b), L has no path from $V(uAv - \{u, v\})$ to $\{b, b'\}$ internally disjoint from $A \cup C$. Let $u', v' \in V(A)$ with $u'Av'$ maximal such that $uAv \subseteq u'Av'$ and L has no path from $V(u'Av' - \{u', v'\})$ to $\{b, b'\}$ internally disjoint from $A \cup C$. Then for any component D^* of L' , $N(D^*) \subseteq u'Av'$ or $N(D^*) \cap V(u'Av' - \{u', v'\}) = \emptyset$. (For otherwise, let $u^*, v^* \in N(D^*) \cap V(A)$ such that a, u^*, u', v^*, v', a' occur on A in this order and $u^*, v^* \notin \{u', v'\}$. Then L has no path from $V(u^*Av' - \{u^*, v'\})$ to $\{b, b'\}$ internally disjoint from $A \cup C$, and so, $\{u^*, v'\}$ contradicts the choice of $\{u', v'\}$.) Also, L^* has no path from $V(u'Av' - \{u', v'\})$ to $V(C)$ internally disjoint from L ; otherwise, by Claim 2, such a path would be from

$aAx_i - x_i$ to $y_{i-1}Cc' - y_{i-1}$ or from $cCy_i - y_i$ to $x_{i-1}Aa' - x_{i-1}$, and internally disjoint from L , a contradiction. Hence $G - \{u', v'\}$ has a component containing no vertex of $V(A \cup C) \cup \{v_0, \dots, v_m\}$, contradicting (f).

Claim 5. No R_i is connected.

Suppose that R_i is connected. By Claim 4, $v_{i-1} = v_i$. By Claim 2, $v_{i-1} \neq b$ or $v_i \neq b'$. By symmetry, assume that $v_{i-1} \neq b$.

If L^* has no path from $aAx_{i-1} - x_{i-1}$ to $y_{i-1}Cc' - y_{i-1}$ internally disjoint from L , and L^* has no path from $cCy_{i-1} - y_{i-1}$ to $x_{i-1}Aa' - x_{i-1}$ internally disjoint from L , then by Claim 3, L^* has a separation (G_1, G_2) such that $V(G_1 \cap G_2) = \{x_{i-1}, v_{i-1}, y_{i-1}\}$, $\{a, b, c\} \subseteq G_1$, and $\{a', b', c'\} \subseteq G_2$, contradicting Claim 1.

So assume (by the symmetry between x_{i-1} and y_{i-1}) that L^* has a path S from $s \in V(aAx_{i-1} - x_{i-1})$ to $s' \in V(y_{i-1}Cc' - y_{i-1})$ internally disjoint from L . Select S so that aAs is minimal. By (1') of (2.3), R_i has a path from y_{i-1} to v_i internally disjoint from $A \cup C$. Hence for all $j < i$, $N(R_j - (A_j \cup C_j \cup \{v_{i-1}\})) \cap V(sAx_{i-1} - s) = \emptyset$; otherwise, $L_{j,i}$ has disjoint paths from a, v_{j-1}, c to c', a', v_i , respectively, contradicting (a). Also, by (c) and by Claim 3, L^* has no path from $cCy_{i-1} - y_{i-1}$ to $sAa' - s$ internally disjoint from L .

If $N(L - (A \cup C)) \cap aAs = \emptyset$ and L^* has no path from aAs to $cCy_{i-1} - y_{i-1}$ internally disjoint from L , then by Claim 3, L^* has a separation (G_1, G_2) such that $V(G_1 \cap G_2) = \{a, v_{i-1}, y_{i-1}\}$, $\{a, b, c\} \subseteq G_1$ and $\{a', b', c'\} \subseteq G_2$, contradicting Claim 1.

So let $r \in V(aAs)$ with rAs minimal such that $r \in N(L - (A \cup C))$ or L^* has a path from r to $cCy_{i-1} - y_{i-1}$ internally disjoint from L .

If $r \in N(L - (A \cup C))$, then let $x'_{i-1} = r$. By (b) (with r as $q \in V(A)$ there), for every component D of L' , $N(D) \cap V(A) \subseteq aAx'_{i-1}$ or $N(D) \cap V(A) \subseteq x'_{i-1}Aa'$. If L^* has a path from r to $r' \in V(cCy_{i-1} - y_{i-1})$, then by (d) (with r, r', s, s' as s, s', t, t' , respectively), there is a vertex $x'_{i-1} \in rAs$ such that, for every component D of L' , $N(D) \cap V(A) \subseteq aAx'_{i-1}$ or $N(D) \cap V(A) \subseteq x'_{i-1}Aa'$.

Hence by the choices of r and s and by Claim 3, L^* has a separation (G_1, G_2) such that $V(G_1 \cap G_2) = \{x'_{i-1}, v_{i-1}, y_{i-1}\}$, $\{a, b, c\} \subseteq G_1$, and $\{a', b', c'\} \subseteq G_2$, contradicting Claim 1.

Claim 6. A_i, C_i and $\{v_{i-1}, v_i\}$ are contained in three different components of R_i .

Suppose on the contrary that Claim 6 fails. By Claim 4 and by symmetry, assume that R_i is not connected, but $\{x_{i-1}, v_{i-1}\} \cup \{x_i, v_i\}$ is contained in a component of R_i . Then $y_{i-1} = y_i$ (since $(R_i, \{x_{i-1}, y_{i-1}\}, \{x_i, y_i\}, (v_{i-1}, v_i))$ is a rung). By (e), $v_{i-1} = v_i$. Hence, $x_{i-1} \neq x_i$, and so, we have $x_{i-1} \neq a'$ and $x_i \neq a$. Moreover, either $\{x_{i-1}, v_{i-1}\} \neq \{a, b\}$ or $\{x_i, v_i\} \neq \{a', b'\}$; otherwise, $b = v_{i-1} = v_i = b'$, contradicting Claim 2. By symmetry, assume that $\{a, b\} \neq \{x_{i-1}, v_{i-1}\} \neq \{a', b'\}$.

If $i = 1$ then define $B = \{v_0\}$, and if $i \geq 2$ then define $B = \bigcup_{1 \leq k \leq i-1} (R_k - (A_k \cup C_k))$. Let $B' = \bigcup_{k=i}^m (R_k - (A_k \cup C_k))$. Then $B \cap B' = \{v_{i-1}\}$, $N(B) \subseteq aAx_{i-1} \cup cCy_{i-1}$, and

$N(B') \subseteq x_{i-1}Aa' \cup y_{i-1}Cc'$.

Define $s \in V(C)$ with sCc' minimal such that $s = c$, or $s \in N(B - v_{i-1})$, or L^* has a path from $aAx_{i-1} - x_{i-1}$ to s internally disjoint from L . Define $t \in V(C)$ with cCt minimal such that $t = c'$, or $t \in N(B' - v_{i-1})$, or L^* contains a path from $x_{i-1}Aa' - x_{i-1}$ to t internally disjoint from L . We distinguish three cases. Note that by (2) of (2.3), $R_i - ((A_i - x_{i-1}) \cup C_i)$ has a path from x_{i-1} to v_i .

Case 1. $s = c$ or $t = c'$.

Suppose $s = c$ (respectively, $t = c'$). Then by Claim 3 and by the choice of s (respectively, t), L^* has a separation (G_1, G_2) such that $V(G_1 \cap G_2) = \{x_{i-1}, v_{i-1}, c\}$ (respectively, $V(G_1 \cap G_2) = \{x_{i-1}, v_{i-1}, c'\}$), $\{a, b, c\} \subseteq G_1$, and $\{a', b', c'\} \subseteq G_2$, contradicting Claim 1.

Case 2. $t \neq c'$ and $s \in N(B - v_{i-1})$.

Then $N(B') \cap V(cCs - s) = \emptyset$ because $N(B') \cap V(C) \subseteq y_{i-1}Cc'$ and $N(B) \cap V(C) \subseteq cCy_{i-1}$. By (b) (with s as $q \in V(C)$), for every component D of L' , $N(D) \cap V(C) \subseteq cCs$ or $N(D) \cap V(C) \subseteq sCc'$. Let $s \in N(R_l - (A_l \cup C_l \cup \{v_l\}))$ for some $l \leq i - 1$.

Hence, L^* does not contain any path from $x_{i-1}Aa' - x_{i-1}$ to $cCs - s$; otherwise, $L_{l,i}$ contains three disjoint paths from a, v_{l-1}, c to $v_{i-1} = v_i, c', a'$, respectively, contradicting (a).

By the choice of s and by Claim 3, L^* has a separation (G_1, G_2) such that $V(G_1 \cap G_2) = \{x_{i-1}, v_{i-1}, s\}$, $\{a, b, c\} \subseteq G_1$, and $\{a', b', c'\} \subseteq G_2$, contradicting Claim 1.

Case 3. $t \neq c'$ and L^* has a path S from $aAx_{i-1} - x_{i-1}$ to s internally disjoint from L .

If L^* has a path from $x_{i-1}Aa' - x_{i-1}$ to t internally disjoint from L , then $t \in sCc'$ by (c) (with s, t as t, t' in (c), respectively). By (d) (with s, t as t, t' in (d), respectively), there is some vertex $q \in sCt$ such that, for every component D of L' , $N(D) \cap V(C) \subseteq cCq$ or $N(D) \cap V(C) \subseteq qCc'$. So by Claim 3 and by the choices of s and t , L^* has a separation (G_1, G_2) such that $V(G_1 \cap G_2) = \{x_{i-1}, v_{i-1}, q\}$, $\{a, b, c\} \subseteq G_1$, and $\{a', b', c'\} \subseteq G_2$, contradicting Claim 1.

Now assume that $t \in N(B' - v_{i-1})$. Then by (b) (with t as $q \in V(C)$), for every component D of L' , $N(D) \cap V(C) \subseteq cCt$ or $N(D) \cap V(C) \subseteq tCc'$. Let $t \in N(R_l - (A_l \cup C_l \cup \{v_{l-1}\}))$ for some l with $i \leq l \leq m$. Hence, $t \in sCc'$; otherwise, $L_{i,l}$ contains three disjoint paths from a, v_{i-1}, c to c', a', v_l , respectively, contradicting (a). Thus, by Claim 3 and by the choices of s and t , L^* has a separation (G_1, G_2) such that $V(G_1 \cap G_2) = \{x_{i-1}, v_{i-1}, t\}$, $\{a, b, c\} \subseteq G_1$, and $\{a', b', c'\} \subseteq G_2$, contradicting Claim 1. This proves Claim 6.

By Claim 6, A, C and $\{v_0, \dots, v_m\}$ are in different components of L . Hence, $A \cup C$ and $\{b, b'\}$ are contained in different components of L^* . Since $b \neq b'$ (by Claim 2), G has a separation (G_1, G_2) such that $V(G_1 \cap G_2) = \{a, b', c\}$, $\{a, b, c\} \subseteq G_1$, and $\{a', b', c'\} \subseteq G_2$, contradicting Claim 1. \square

3 The proof of main result

First, we prove the following lemma, which proves the necessary part of (1.3).

(3.1) Lemma. *Let $(G, \{a, c\}, \{a', c'\}, (b, b'))$ be an obstruction. Then one of the following statements holds.*

- (1) G has a separation (G_1, G_2) of order at most 2 such that $\{a, b, c\} \subseteq V(G_1)$ and $\{a', b', c'\} \subseteq V(G_2)$.
- (2) There is a subset $T \subseteq V(G)$ such that $|T| \leq 3$ and some component of $G - T$ contains no element of $\{a, b, c\} \cup \{a', b', c'\}$.
- (3) $(G, (a, b, c), (a', b', c'))$ is a ladder.
- (4) G has a separation (J, L) such that $V(J \cap L) = \{w_0, \dots, w_n\}$, (J, w_0, \dots, w_n) is 3-planar, $(L, (a, b, c), (a', b', c'))$ is a ladder along a sequence $v_0 \dots v_m$, where $v_0 = b$, $v_m = b'$, and $w_0 \dots w_n$ is the reduced sequence of $v_0 \dots v_m$.

Proof. Assume that (1) and (2) do not hold. Let A, B, C be disjoint paths from $\{a, b, c\}$ to $\{a', b', c'\}$ in G such that $a \in A$, $c \in C$, and B is from b to b' . We choose A, B and C to be induced paths. Let H' denote the component of $G - (A \cup C)$ containing B . Let K be the subgraph of G obtained from $A \cup C \cup H'$ by adding those edges in G between $V(A) \cup V(C)$ and $V(H')$. It is straightforward to verify that A, B, C, H' and K satisfy the conditions (i)–(iii) of (2.4) (with K as G there). By applying (2.4), one of (1)–(4) of (2.4) holds.

Because of A, B, C , (1) of (2.4) does not hold. Since we assume that (2) of (3.1) does not hold, if there is some $T \subseteq V(K)$ with $|T| \leq 3$ such that some component U of $G - T$ contains no element of $\{a, b, c\} \cup \{a', b', c'\}$, then T would be contained in the “ladder part” of K when (3) or (4) of (2.4) holds (see the comment after (1.3)). So we may assume that either (1a) $(K, (a, b, c), (a', b', c'), A, C)$ (when $a' \in A$ and $c' \in C$) or $(K, (a, b, c), (c', b', a'), A, C)$ (when $a' \in C$ and $c' \in A$) is a good ladder along a sequence $v_0 \dots v_m$, where $v_0 = b$ and $v_m = b'$ (in this case, let $J = \{w_0, \dots, w_n\}$ and $L = K$, where $w_0 \dots w_n$ is the reduced sequence of $v_0 \dots v_m$), or (1b) K has a separation (J, L) such that $V(J \cap L) = \{w_0, \dots, w_n\}$, (J, w_0, \dots, w_n) is 3-planar, $(L, (a, b, c), (a', b', c'), A, C)$ (when $a' \in A$ and $c' \in C$) or $(L, (a, b, c), (c', b', a'), A, C)$ (when $a' \in C$ and $c' \in A$) is a good ladder along $v_0 \dots v_m$, where $v_0 = b$, $v_m = b'$, and $w_0 \dots w_n$ is the reduced sequence of $v_0 \dots v_m$.

Let $(R_i, (x_{i-1}, v_{i-1}, y_{i-1}), (x_i, v_i, y_i), A_i, C_i)$, $i = 1, \dots, m$, be the good rungs of L , where $A_i = A \cap R_i$, $C_i = C \cap R_i$, $a = x_0$, $c = y_0$, $a' = x_m$ (when $a' \in A$) or $c' = x_m$ (when $c' \in A$), and $c' = y_m$ (when $c' \in C$) or $a' = y_m$ (when $a' \in C$). Let $v(L) = \{x \in V(R_k - v_k) : v_{k-1} = v_k, 1 \leq k \leq m\}$. We select A, B, C, H', J and L such that

- (I) H' is maximal,

(II) subject to (I), $v(L)$ is maximal, and

(III) subject to (II), J is maximal.

Without loss of generality, we may choose the notation so that A, B, C are from a, b, c to a', b', c' , respectively. Let $L' = G - K$ and let $L^* = G[V(L \cup L')]$. Then L^* is obtained from the disjoint union of L and L' by adding edges between $V(A) \cup V(C)$ and $V(L')$ or between $V(A)$ and $V(C)$. Next we will show that L^*, L, L', A, C satisfy the remaining conditions (a)–(f) of (2.5).

Note that since (2) does not hold for G , J can be drawn in a closed disc in the plane with no edge crossings such that w_0, w_1, \dots, w_n occur on the boundary of the disc in that cyclic order.

- (a) For any $i, j \in \{1, \dots, m\}$ with $i \leq j$, let $L_{i,j}$ denote the graph obtained from L^* by deleting all $R_k - (A_k \cup C_k \cup \{v_{i-1}, v_j\})$ with $k < i$ or $k > j$ and by deleting v_0, \dots, v_m except v_{i-1} and v_j . Then $(L_{i,j}, \{a, c\}, \{a', c'\}, (v_{i-1}, v_j))$ is an obstruction.

Clearly (a) holds when $v_{i-1} = v_j$. So assume that $v_{i-1} \neq v_j$. For convenience, let G^* denote the union of J and all $R_k - (A_k \cup C_k)$ with $k < i$ or $k > j$. Then $V(G^* \cap L_{i,j}) = \{v_{i-1}, v_j\}$.

If we let $\mathcal{J} = \{V(R_k) - (V(A_k \cup C_k) \cup \{v_{k-1}, v_k\}) : k < i \text{ or } k > j\}$, then $(G^*, \mathcal{J}, v_0, \dots, v_{i-1}, v_j, \dots, v_m)$ is 3-planar. Hence, let \mathcal{J}^* be a collection of pairwise disjoint subsets of $V(J) - \{v_0, \dots, v_{i-1}, v_j, \dots, v_m\}$ such that $(G^*, \mathcal{J}^*, v_0, \dots, v_{i-1}, v_j, \dots, v_m)$ is 3-planar and \mathcal{J}^* is minimal.

If G^* contains a $(\{b, v_{i-1}\}, \{b', v_j\})$ -linkage, then $(L_{i,j}, \{a, c\}, \{a', c'\}, (v_{i-1}, v_j))$ is an obstruction. For otherwise, $L_{i,j}$ has three disjoint paths from $\{a, v_{i-1}, c\}$ to $\{a', v_j, c'\}$, none from v_{i-1} to v_j . These three paths in $L_{i,j}$ combined with a $(\{b, v_{i-1}\}, \{b', v_j\})$ -linkage in G^* give three disjoint paths in G from $\{a, b, c\}$ to $\{a', b', c'\}$, none from b to b' . This contradicts the assumption that $(G, \{a, c\}, \{a', c'\}, (b, b'))$ is an obstruction.

So assume that G^* contains no $(\{b, v_{i-1}\}, \{b', v_j\})$ -linkage. Then $p(G^*, \mathcal{J}^*)$ has no $(\{b, v_{i-1}\}, \{b', v_j\})$ -linkage; otherwise, by (2.2) (with $G^*, \mathcal{J}^*, v_{i-1}, v_j$ as G, \mathcal{A}, u, v , respectively), G^* would contain a $(\{b, v_{i-1}\}, \{b', v_j\})$ -linkage. By planarity, $p(G^*, \mathcal{J}^*)$ has a separation (G'_1, G'_2) such that $|V(G'_1 \cap G'_2)| \leq 1$, $\{v_{i-1}, v_j\} \subseteq G'_1$, and $\{b, b'\} \subseteq G'_2$. Again by planarity of $p(G^*, \mathcal{J}^*)$, there are $r \leq i-1$ and $s \geq j$ such that $\{v_r, \dots, v_s\} \subseteq G'_1$ and $\{v_0, \dots, v_{r-1}, v_{s+1}, \dots, v_m\} \subseteq G'_2$. Let G_i denote the subgraph of G^* induced by the vertices contained in $V(G'_i)$ or contained in some $U \in \mathcal{J}^*$ with $N_{G^*}(U) \subseteq V(G'_i)$. Then (G_1, G_2) is a separation of G^* such that $|V(G_1 \cap G_2)| \leq 1$, $\{v_r, \dots, v_s\} \subseteq G_1$, and $\{v_0, \dots, v_{r-1}, v_{s+1}, \dots, v_m\} \subseteq G_2$.

Because $K - (A \cup C)$ is connected, $|V(G_1 \cap G_2)| = 1$. Let w be the vertex in $V(G_1 \cap G_2)$. Then K has a separation (M_1, M_2) such that $V(M_1 \cap M_2) = \{x_{r-1}, y_{r-1}, w, x_s, y_s\}$, $M_1 = G_1 \cup (\bigcup_{r \leq k \leq s} R_k)$, and $G_2 \subseteq M_2$.

Let a_0, \dots, a_l be vertices in that order on A and c_0, \dots, c_l be vertices in that order on C such that (a1) $a_0 = x_{r-1}$, $a_l = x_s$, $c_0 = y_{r-1}$, and $c_l = y_s$, (a2) for $i = 1, \dots, l-1$, $\{a_i, c_i\}$ is a 2-cut of $M_1 - w$, and (a3) for any $x \in V(a_{i-1}Aa_i)$ and $y \in V(c_{i-1}Cc_i)$, $\{x, y\}$ is not a 2-cut of $M_1 - w$ unless $\{x, y\} = \{a_{i-1}, c_{i-1}\}$ or $\{x, y\} = \{a_i, c_i\}$. Let S denote the set of edges of M_1 with both ends contained in $\{a_k, c_k, w\}$, $k = 0, \dots, l$.

Let M_1^k denote the subgraph of $M_1 - S$ induced by those vertices x such that every path in $M_1 - w$ from x to $\{x_{r-1}, y_{r-1}, w, x_s, y_s\}$ intersect $\{a_{k-1}, a_k, w, c_{k-1}, c_k\}$. Because (2) does not hold for G and by (a3), M_1^k has no 3-separation (M', M'') such that $\{a_{k-1}, w, c_{k-1}\} \subseteq M'$ and $\{a_k, w, c_k\} \subseteq M''$. By (1) of (1.1), $(M_1^k, (a_{k-1}, w, c_{k-1}), (a_k, w, c_k))$ is a rung. In fact, by (a3), it is easy to show that $(M_1^k, (a_{k-1}, w, c_{k-1}), (a_k, w, c_k), a_{k-1}Aa_k, c_{k-1}Cc_k)$ is a good rung.

Again because (2) does not hold for G and by (a3), $M_1 = (\bigcup_{k=1}^l M_1^k) + S$. Therefore, we see that $(M_1, (x_{r-1}, y_{r-1}), (x_s, y_s), A \cap M_1, C \cap M_1)$ is a good ladder along w .

Let $L_1 = M_1 \cup L$, and let $J_1 = J - (G_1 - w)$. Then (J_1, L_1) is a separation of K such that $V(J_1 \cap L_1) = \{w_0, \dots, w_{t-1}, w, w_t, \dots, w_n\}$ for some appropriate t , $(J_1, w_0, \dots, w_{t-1}, w, w_t, \dots, w_n)$ is 3-planar, $(L_1, (a, b, c), (a', b', c'), A, C)$ is a good ladder along a sequence whose reduced sequence is $w_0 \dots w_{t-1} w w_t \dots w_n$. Note that H' remains unchanged, but $v(L) \cup (\{v_{i-1}, v_j\} - \{w\}) \subseteq v(L_1)$, and $v_{i-1}, v_j \notin v(L)$, contradicting (II).

- (b) If $q \in V(A)$ (respectively, $q \in V(C)$) such that L contains a path from q to $\{v_0, \dots, v_m\}$ internally disjoint from $V(A \cup C) \cup \{v_0, \dots, v_m\}$, then for every component D of L' , $N(D) \cap V(A) \subseteq aAq$ or $N(D) \cap V(A) \subseteq qAa'$ (respectively, $N(D) \cap V(C) \subseteq cCq$ or $N(D) \cap V(C) \subseteq qCc'$).

Suppose (b) is false. By symmetry, assume that L contains a path from $q \in V(A)$ to $\{v_0, \dots, v_m\}$ internally disjoint from $V(A \cup C) \cup \{v_0, \dots, v_m\}$, and assume that there is a component D of L' such that $N(D) \cap V(aAq - q) \neq \emptyset$ and $N(D) \cap V(qAa' - q) \neq \emptyset$. Let $u \in N(D) \cap V(aAq - q)$ and $v \in N(D) \cap V(qAa' - q)$.

Let A' be an induced path from a to a' in L^* such that $V(A') \subseteq aAu \cup D \cup vAa'$. Then the component of $G - (A' \cup C)$ containing B is larger than H' , contradicting (I).

- (c) If L^* contains paths from distinct $s, s' \in V(A)$ (in that order from a to a') to distinct $t, t' \in V(C)$ (in that order from c' to c), respectively, which are internally disjoint from L , then L contains no path from $V(sAs' - \{s, s'\}) \cup V(tCt' - \{t, t'\})$ to $\{v_0, \dots, v_m\}$ internally disjoint from $V(A \cup C) \cup \{v_0, \dots, v_m\}$.

Let P, P' be the paths in (c) from s, s' to t, t' , respectively, which are internally disjoint from L . If $P \cap P' \neq \emptyset$, then there is a component D of L' such that $\{s, s', t, t'\} \subseteq N(D)$, and (c) follows from (b). So assume that $P \cap P' = \emptyset$.

Suppose that (c) fails, and assume by symmetry that L^* has a path from $V(sAs' - \{s, s'\})$ to $\{v_0, \dots, v_m\}$ internally disjoint from $V(A \cup C) \cup \{v_0, \dots, v_m\}$. Let A' be an

induced path in L^* from a to c' such that $V(A') \subseteq aAs \cup P \cup tCc'$, and let C' be an induced path in L^* from c to a' such that $V(C') \subseteq cCt' \cup P' \cup s'Aa'$. Then the component of $G - (A' \cup C')$ containing B is larger than H' , contradicting (I).

- (d) Suppose that L^* contains paths from $s, s' \in V(A)$ (in that order from a to a') to $t, t' \in V(C)$ (in that order from c to c') which are internally disjoint from L . If L contains a path from $V(tCt' - \{t, t'\})$ (respectively, $V(sAs' - \{s, s'\})$) to $\{v_0, \dots, v_m\}$ internally disjoint from $V(A \cup C) \cup \{v_0, \dots, v_m\}$, then there is a vertex $q \in sAs'$ (respectively, $q \in tCt'$) such that, for every component D of L' , $N(D) \cap V(A) \subseteq aAq$ or $N(D) \cap V(A) \subseteq qAa'$ (respectively, $N(D) \cap V(C) \subseteq cCq$ or $N(D) \cap V(C) \subseteq qCc'$).

By symmetry, assume L contains a path from $V(tCt' - \{t, t'\})$ to $\{v_0, \dots, v_m\}$ internally disjoint from $V(A \cup C) \cup \{v_0, \dots, v_m\}$. Let $r, r' \in V(A)$ with $rAr' \subseteq sAs'$ and rAr' minimal such that L^* has paths R, R' from $r, r' \in V(A)$ to t, t' , respectively, which are internally disjoint from L . By (c), a, s, r, r', s', a' occur on A in that order. We will show that there is some $q \in V(rAr')$ such that, for every component D of L' , $N(D) \cap V(A) \subseteq aAq$ or $N(D) \cap V(A) \subseteq qAa'$.

We may assume that L' has a component D_1 such that $N(D_1) \cap V(aAr - r) \neq \emptyset$ and $N(D_1) \cap V(rAa' - r) \neq \emptyset$; otherwise, $q := r$ would be the desired vertex. Let $p_1 \in V(aAr - r) \cap N(D_1)$ and $q_1 \in V(rAa' - r) \cap N(D_1)$, and let P_1 be a path in L^* from p_1 to q_1 through D_1 and internally disjoint from L . Select D_1, P_1, p_1, q_1 so that q_1Aa' is minimal and subject to this, p_1Aq_1 is maximal. See Figure 1.

Then $q_1 \in rAr' - r$. For otherwise, $q_1 \in r'Aa' - r'$. By (c), $D_1 \cap (R \cup R') = \emptyset$. Let A' be an induced path in L^* from a to a' such that $V(A') \subseteq aAp_1 \cup P_1 \cup q_1Aa'$, and let C' be an induced path from c to c' in L^* such that $V(C') \subseteq cCt \cup R \cup rAr' \cup R' \cup t'Cc'$. Then the component of $G - (A' \cup C')$ containing B is larger than H' , contradicting (I).

By (c), $R' \cap D_1 = \emptyset$. By the minimality of rAr' , $R \cap D_1 = \emptyset$. We may assume that there are $p_2 \in V(aAq_1 - q_1)$ and $q_2 \in V(q_1Aa' - q_1)$ such that $\{p_2, q_2\} \subseteq N(D_2)$ for some component D_2 of L' ; otherwise, $q := q_1$ would be the desired vertex. Let P_2 be a path in L^* from p_2 to q_2 through D_2 and internally disjoint from L . Select D_2, P_2, p_2, q_2 such that q_2Aa' is minimal, and subject to this, p_2Aq_2 is maximal. By the choice of D_1, P_1, p_1 and q_1 , we conclude that $p_2 \in rAq_1 - q_1$. If $q_2 \in r'Aa' - r'$, then we stop this process; and if $q_2 \in q_1Ar' - q_1$, then we find D_3, P_3, p_3, q_3 in a similar way. Since G is a finite graph, we have a sequence P_1, P_2, \dots, P_l of paths with P_i from p_i to q_i such that $q_i \in r'Aa' - r'$, and for $i \in \{1, \dots, l-1\}$, $p_{i+1} \in q_{i-1}Aq_i - q_i$ and $q_i \in p_{i+1}Ap_{i+2} - p_{i+1}$, where $q_0 = r$ and $p_{l+1} = r'$. Figure 1 illustrates the cases for $l = 3$ and $l = 4$.

By (c) and by the minimality of rAr' , $(P_i - \{p_i, q_i\}) \cap (R \cup R') = \emptyset$ for $i \in \{1, \dots, l\}$. By the choice of P_i , all P_i 's are internally disjoint, and $P_i \cap P_j = \emptyset$ if $|i - j| \neq 2$.

If l is odd, then let A' be an induced path in L^* from a to a' such that $V(A') \subseteq aAp_1 \cup P_1 \cup q_1Ap_3 \cup P_3 \cup q_3Ap_5 \cup \dots \cup P_l \cup q_lAa'$, and let C' be an induced path in L^* from c

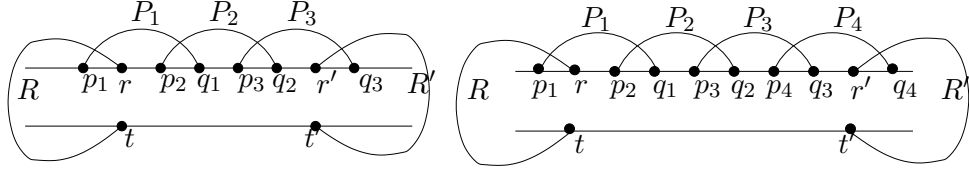


Figure 1: Paths P_1, \dots, P_l .

to c' such that $V(C') \subseteq cCt \cup R \cup rAp_2 \cup P_2 \cup q_2Ap_4 \cup P_4 \cup \dots \cup P_{l-1} \cup q_{l-1}Ar' \cup R' \cup t'Cc'$. If l is even, then let A' be an induced path in L^* from a to c' such that $V(A') \subseteq aAp_1 \cup P_1 \cup q_1Ap_3 \cup P_3 \cup q_3Ap_5 \cup \dots \cup P_{l-1} \cup q_{l-1}Ar' \cup R' \cup t'Cc'$, and let C' be an induced path in L^* from c to a' such that $V(C') \subseteq cCt \cup R \cup rAp_2 \cup P_2 \cup q_2Ap_4 \cup P_4 \cup \dots \cup P_l \cup q_lAa'$. Clearly, the component of $G - (A' \cup C')$ containing B is larger than H' , contradicting (I).

- (e) If R_i is not connected and $\{x_{i-1}, v_{i-1}\} \cup \{x_i, v_i\}$ (respectively, $\{y_{i-1}, v_{i-1}\} \cup \{y_i, v_i\}$) is contained in a component of R_i , then $v_{i-1} = v_i$.

Suppose $v_{i-1} \neq v_i$. Assume by symmetry that $\{x_{i-1}, v_{i-1}\} \cup \{x_i, v_i\}$ is contained in a component U of R_i . Then $y_{i-1} = y_i$ and $|V(U)| \geq 4$ because $E(U) \neq \emptyset$ and R_i has no 3-separation (R', R'') such that $\{x_{i-1}, v_{i-1}, y_{i-1}\} \subseteq R'$ and $\{x_i, v_i, y_i\} \subseteq R''$. By assumption that (2) does not hold for G , $x_{i-1} \neq x_i$;

Moreover, there is a collection \mathcal{U} of pairwise disjoint subsets of $V(U) - \{x_{i-1}, v_{i-1}, x_i, v_i\}$ such that $(U, \mathcal{U}, x_{i-1}, x_i, v_i, v_{i-1})$ is 3-planar; otherwise, U contains an $(\{x_{i-1}, v_i\}, \{v_{i-1}, x_i\})$ -linkage (by (2.1)), and so, $L_{i,i}$ has disjoint paths from a, v_{i-1}, c to v_i, a', c' , respectively, contradicting (a).

Note that $U - A_i$ is connected. We claim that $\{v_{i-1}, v_i\}$ is contained in a block M of $U - A_i$. For otherwise, let (U_1, U_2) be a 1-separation in $U - A_i$ such that $v_{i-1} \in U_1$ and $v_i \in U_2$. Then there are $u_j \in V(A_i) \cap N(U_j)$, $j = 1, 2$, such that x_{i-1}, u_2, u_1, x_i occur on A in that order; as otherwise U has a 2-separation (U'_1, U'_2) such that $U_1 \cup \{x_{i-1}\} \subseteq U'_1$ and $U_2 \cup \{x_i\} \subseteq U'_2$, and $(U'_1 \cup \{y_i\}, U'_2 \cup \{y_i\})$ would be a 3-separation in R_i , a contradiction. Hence, U has an $(\{x_{i-1}, v_i\}, \{v_{i-1}, x_i\})$ -linkage. This implies that $L_{i,i}$ has disjoint paths from a, v_{i-1}, c to v_i, a', c' , respectively, contradicting (a).

We claim that $X \cap V(M) = \emptyset$ for all $X \in \mathcal{U}$. Otherwise, let $X \in \mathcal{U}$ such that $X \cap V(M) \neq \emptyset$. Then $|N_U(X) \cap V(M)| \geq 2$ because M is 2-connected. Hence, $|N_U(X) \cap V(A)| \leq 1$. Therefore $G - N_U(X)$ has a component containing no element of $\{a, b, c\} \cup \{a', b', c'\}$, contradicting the assumption that (2) does not hold for G .

Thus M is a plane subgraph of $p(U, \mathcal{U})$. Let b_0, \dots, b_k be the vertices on the outer cycle of M (in that clockwise order from $b_0 = v_{i-1}$ to $b_k = v_i$) such that, for each $j \in \{1, \dots, k-1\}$, $b_j \in N(A)$ or b_j is a cut vertex of $U - A$. Let $R'_i = R_i - (M - \{b_0, \dots, b_k\})$.

By 3-planarity of U , it is easy to verify that $(R'_i, (x_{i-1}, v_{i-1}, y_i), (x_i, v_i, y_i))$ is a ladder along a sequence whose reduced sequence is $b_0 \dots b_k$. In fact, it is easy to see that $(R'_i, (x_{i-1}, v_{i-1}, y_i), (x_i, v_i, y_i), A_i, C_i)$ is a good ladder (its rungs are not connected).

Let $J_1 = J \cup M$ and let $L_1 = L - (M - \{b_0, \dots, b_k\})$. Then (J_1, L_1) is a separation in K such that $V(L_1 \cap J_1) = \{w_0, \dots, w_n\} \cup \{b_0, \dots, b_k\}$, $(L_1, w_0, \dots, w_{j-1}, b_0, \dots, b_k, w_j, \dots, w_n)$ is 3-planar (for some appropriate j such that $w_{j-1} = v_r$ and $w_j = v_s$ for some $r \leq i-1$ and $s \geq i$), $(L_1, (a, b, c), (a', b', c'), A, C)$ is a good ladder along a sequence S , where $w_0 \dots w_{j-1} b_0 \dots b_k w_j \dots w_n$ is the reduced sequence of S . Note that H' remains the same and $v(L) \subseteq v(L_1)$ (since $v_{i-1} \neq v_i$), but J_1 contains J as a proper subgraph, contradicting (III). This proves (e).

Because (2) does not hold for G , we have

- (f) For any $T \subseteq V(L^*)$ with $|T| \leq 3$, every component of $L^* - T$ contains a vertex in $V(A \cup C) \cup \{v_0, \dots, v_m\}$.

By (a)–(f), L^*, L, L', A, C satisfy all conditions of (2.5). Hence, by (2.5), $(L^*, (a, b, c), (a', b', c'))$ is a ladder along a sequence $z_0 z_1 \dots z_p$ whose reduced sequence is the reduced sequence of $v_0 \dots v_m$. Hence $w_0 \dots w_n$ is the reduced sequence of $z_0 \dots z_p$. Therefore, either $G = L^*$ or G has a separation (J, L^*) such that $V(J \cap L^*) = \{w_0, \dots, w_n\}$, (J, w_0, \dots, w_n) is 3-planar, and $(L^*, (a, b, c), (a', b', c'))$ is a ladder along $z_0 \dots z_p$, where $z_0 = b$, $z_p = b'$, $w_0 \dots w_n$ is the reduced sequence of $z_0 \dots z_p$. \square

Proof of (1.3). Suppose $(G, \{a, c\}, \{a', c'\}, (b, b'))$ is an obstruction, and assume that, for any $T \subseteq V(G)$ with $|T| \leq 3$, every component of $G - T$ contains an element of $\{a, b, c\} \cup \{a', b', c'\}$. Then by (3.1), (1), (2) or (3) of (1.3) holds.

Now assume that (1), (2) or (3) of (1.3) holds. We wish to show that $(G, \{a, c\}, \{a', c'\}, (b, b'))$ is an obstruction.

If (1) of (1.3) holds, then $(G, \{a, c\}, \{a', c'\}, (b, b'))$ is an obstruction in the trivial sense. If (2) of (1.3) holds, then $(G, (a, b, c), (a', b', c'))$ is a ladder, and so, an obstruction. So we may assume that (1) and (2) of (1.3) do not hold, and (3) of (1.3) holds. Then G contains three disjoint paths from $\{a, b, c\}$ to $\{a', b', c'\}$. Let A, B, C be disjoint paths in G from $\{a, b, c\}$ to $\{a', b', c'\}$, with $b \in B$. We may assume that A, B, C are chosen so that $b' \notin B$; for otherwise, $(G, (a, b, c), (a', b', c'))$ is an obstruction.

Let (J, L) be a separation of G such that $V(J \cap L) = \{w_0, \dots, w_n\}$, (J, w_0, \dots, w_n) is 3-planar, and $(L, (a, b, c), (a', b', c'))$ is a ladder with rungs $(R_i, (x_{i-1}, v_{i-1}, y_{i-1}), (x_i, v_i, y_i))$, $i = 1, \dots, m$, where $x_0 = a$, $v_0 = b$, $y_0 = c$, $x_m = a'$, $v_m = b'$, $y_m = c'$, and $w_0 \dots w_n$ is the reduced sequence of $v_0 \dots v_m$. We use induction on the number of rungs in L .

Since $b' \notin B$, $B \not\subseteq J$. Hence, let $w_k \in B$ with k maximum such that $bBw_k \subseteq J$. Then $k < n$; otherwise, $b' \in B$ (since (J, w_0, \dots, w_n) is 3-planar). Let $v_s = w_k$ with s minimum and let $v_t = w_{k+1}$ with t maximum. Let L^* be the union of $\bigcup_{s < i \leq t} R_i$ and those edges of L with both ends in some $\{x_i, v_i, y_i\}$, $s < i \leq t$. Then $(L^*, (x_s, v_s, y_s), (x_t, v_t, y_t))$ is

a ladder along $v_s \dots v_t$. Since every path in G from $\{a, b, c\}$ to $\{a', b', c'\}$ must intersect $V(bBw_k) \cup \{x_s, y_s\}$, $A \cap L^*$, $B \cap L^*$, and $C \cap L^*$ are disjoint paths from $\{x_s, w_k, y_s\}$ to $\{x_t, w_{k+1}, y_t\}$. Since $(L^*, (x_s, v_s, y_s), (x_t, v_t, y_t))$ is a ladder, $B \cap L^*$ is a path from w_k to w_{k+1} .

Let L' denote the union of $\bigcup_{j=t+1}^m R_j$ and edges in L with both ends in some $\{x_j, v_j, y_j\}$, $j = t+1, \dots, m$. Then $(L', (x_t, v_t, y_t), (a', b', c'))$ is a ladder along $v_t \dots v_m$. Let $J' = J + w_k w_{k+1}$, and let the edge $w_k w_{k+1}$ be added such that (J', w_0, \dots, w_n) is 3-planar. It is easy to see that $A \cap (J' \cup L')$, $(B \cap (J' \cup L')) + w_k w_{k+1}$, $C \cap (J' \cup L')$ are disjoint paths in $J' \cup L'$ from $\{b, x_t, y_t\}$ to $\{a', b', c'\}$. Note that (J', L') is a separation in $J' \cup L'$ such that $V(J' \cap L') = \{b, w_{k+1}, \dots, w_n\}$, $(J', b, w_{k+1}, \dots, w_n)$ is 3-planar, and $(L', (x_t, b, y_t), (a', b', c'))$ is a ladder along $bv_t \dots v_m$, where $v_m = w_n = b'$, and $bw_{k+1} \dots w_n$ is the reduced sequence of $bv_t \dots v_m$. Since L' has fewer rungs than L , it follows from induction that $(B \cap (J' \cup L')) + w_k w_{k+1}$ is from b to b' . Hence, B is from b to b' . So $(G, \{a, c\}, \{a', c'\}, (b, b'))$ is an obstruction. \square

4 Connectivity

In this section, we prove (1.4), and give an example to show that (1.4) is best possible.

Proof of (1.4). Suppose that G is 8-connected and $\{a, b, c, a', b', c'\} \subseteq V(G)$, and assume that $(G, \{a, c\}, \{a', c'\}, (b, b'))$ is an obstruction. By (1.3), either (i) $(G, (a, b, c), (a', b', c'))$ is a ladder along a sequence $v_0 \dots v_m$ (in this case, let $J = \{w_0, \dots, w_n\}$ and $L = G$, where $w_0 \dots w_n$ is the reduced sequence of $v_0 \dots v_m$), or (ii) G has a separation (J, L) such that $V(J \cap L) = \{w_0, \dots, w_n\}$, (J, w_0, \dots, w_n) is 3-planar, and $(L, (a, b, c), (a', b', c'))$ is a ladder with rungs along $v_0 \dots v_m$, where $w_0 \dots w_n$ is the reduced sequence of $v_0 \dots v_m$.

Then $L - \{w_0, \dots, w_n\} \neq \emptyset$; otherwise, $G = J$ is 3-planar, and so, is at most 5-connected, a contradiction. Since (J, w_0, \dots, w_n) is 3-planar and since G is 8-connected, J is a plane graph. Let $(R_i, (x_{i-1}, v_{i-1}, y_{i-1}), (x_i, v_i, y_i))$, $i = 1, \dots, m$, be the rungs of L , where $x_0 = a$, $v_0 = b$, $y_0 = c$, $x_m = a'$, $v_m = b'$, and $y_m = c'$.

(a) We claim that L is at most 6-connected.

This is obvious when $m \geq 2$. So assume that $m = 1$. Then $(L, (a, b, c), (a', b', c'))$ is a rung. Since $\{a, c\} \neq \{a', c'\}$ and $b \neq b'$, L is at most 6-connected.

By (a) and since G is 8-connected, $J - \{w_0, \dots, w_n\} \neq \emptyset$, $n \geq 7$, and $V(R_i) = \{x_{i-1}, y_{i-1}, v_{i-1}\} \cup \{x_i, y_i, v_i\}$ for $i \in \{1, \dots, m\}$.

(b) We claim that $v_0 \neq v_1$ and $\{x_1, y_1\} = \{x_0, y_0\}$, and $v_{m-1} \neq v_m$ and $\{x_m, y_m\} = \{x_{m-1}, y_{m-1}\}$.

If $\{x_1, y_1\} \neq \{x_0, y_0\}$, then $\{v_0, x_1, v_1, y_1\}$ is a cut set of G , a contradiction. So $\{x_1, y_1\} = \{x_0, y_0\}$, and hence, $v_1 \neq v_0$.

Similarly, $v_{m-1} \neq v_m$ and $\{x_m, y_m\} = \{x_{m-1}, y_{m-1}\}$.

(c) We claim that, for every $x \in V(J) - \{w_1, \dots, w_{n-1}\}$, x is adjacent to at most two vertices in $\{w_1, \dots, w_{n-1}\}$.

Suppose that x is adjacent to at least three vertices of $\{w_1, \dots, w_{n-1}\}$. Let $i, j \in \{1, \dots, n-1\}$ with $i \leq j-2$ such that w_i, w_j are neighbors of x . Let $v_k = w_i$ and $v_l = w_j$. Then $\{v_k, x_k, y_k, v_l, x_l, y_l, x\}$ is a cut set of G , a contradiction.

Now, consider $J' = J - \{w_1, \dots, w_{n-1}\}$. Since J is a plane graph, J' is a plane graph. Let $x \in V(J')$. If $x \in \{w_0, w_n\}$, then by (b) and (c) (since $w_0 = v_0$ and $w_n = v_m$), the degree of x in J' is at least 3. If $x \notin \{w_0, w_n\}$, then by (c), the degree of x in J' is at least 6. Thus the number of edges in J' is at least $\frac{6(|V(J')|-2)+2 \cdot 3}{2} = 3|V(J')| - 3$, contradicting the planarity of J' . \square

Next, we show that (1.4) is best possible by constructing 7-connected obstructions. A *near triangulation* is a plane graph in which every face, with the possible exception of the infinite face, is bounded by a triangle.

First, we construct a near triangulation J whose outer cycle C contains an edge b_1b_2 (with $E(b_1Cb_2) = \{b_1b_2\}$) such that

- (1) $d(x) \geq 3$ for all $x \in V(J)$,
- (2) $d(x) \geq 7$ for all $x \in V(J - C) \cup \{b_1, b_2\}$,
- (3) if T is a cut set of J with $|T| \leq 6$, then $|T| \geq 3$ and every component of $J - T$ contains a vertex of C , and
- (4) for each $x \in V(C) - \{b_1, b_2\}$, there does not exist $T \subseteq V(J)$ with $|T| \leq 4$ such that every path from b_2Cx to xCb_1 in $J - b_1b_2$ intersects T .

We start with a path $J_0 = P_0$ consisting of a single edge b_1b_2 . Suppose we have constructed a near triangulation J_i and a path P_i in J_i for some $i \geq 0$ such that $P_0 \cap P_i = \emptyset$ and $P_0 \cup P_i$ is contained in the outer cycle of J_i . We construct a near triangulation L_{i+1} as follows. (See Figure 2 for an illustration with $i = 3$.)

Let P_{i+1} be a sufficiently long path (for example, at least 6 times as long as P_i). In the disjoint union of P_{i+1} and L_i , we add edges from each vertex of P_i to at least six consecutive vertices on P_{i+1} so that $P_{i+1} \cup P_0$ is contained in the outer cycle of J_{i+1} , $d(x) \geq 7$ if $x \in V(J_i)$, $d(x) \geq 3$ if $x \in V(P_{i+1})$ and x is not an end of P_{i+1} , and $d(x) = 2$ if x is an end of P_{i+1} .

Let x and x' be the ends of P_4 . Add a sufficiently long path Q from x to x' such that Q is internally disjoint from J_4 . We add edges from each vertex of $P_4 - \{x, x'\}$ to at least six consecutive vertices on Q such that the result is a near triangulation J , $Q \cup P_0$ is contained in its outer cycle, and $d(x) \geq 7$ if $x \in V(J_4)$ and $d(x) \geq 3$ if $x \in V(Q)$. It is

straightforward to check that the resulting graph J satisfying (1)–(4) above. We relabel the vertices in J so that $C = w_0w_1 \dots w_nw_0$ is the outer cycle of J , where $w_0 = b_2 = b$ and $w_n = b_1 = b'$. See Figure 2. We may choose n to be odd.

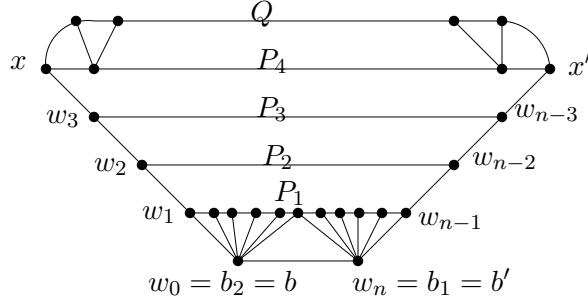


Figure 2: A near triangulation.

Next, we construct a ladder along $v_0 \dots v_m$ whose reduced sequence is $w_0 \dots w_n$, where $m = 2n - 3$ and $n \geq 6$ is an odd integer, and describe how it is attached to J along $w_0 \dots w_n$. See Figure 3.

Let $v_i = w_i$ for $i = 0, 1, 2$, and let $v_{m-2} = w_{n-2}$, $v_{m-1} = w_{n-1}$, and $v_m = w_n$. For $i = 1, 2, m - 1, m$, let $x_{i-1} = x_i$ and $y_{i-1} = y_i$, and let R'_i be the complete graph on $\{x_i, y_i, v_{i-1}, v_i\}$. For each $i = 2k + 1$, $3 \leq i \leq m - 2$, let $x_{i-1} \neq x_i$, $y_{i-1} \neq y_i$, and $v_{i-1} = v_i = w_{k+1}$, and let R'_i be the complete graph on $\{x_{i-1}, x_i, v_i, y_{i-1}, y_i\}$. For each $i = 2k + 2$ with $4 \leq i \leq m - 3$, let $x_{i-1} = x_i$, $y_{i-1} = y_i$, $v_{i-1} = w_{k+1}$ and $v_i = w_{k+2}$, and let R'_i be the complete graph on $\{x_i, v_{i-1}, v_i, y_i\}$. Let S be the set of 2-element subsets of $\{x_i, v_i, y_i\}$ for all $i = 0, \dots, m$. Let $R_i = R'_i - S$. Then $(R_i, (x_{i-1}, v_{i-1}, y_{i-1}), (x_i, v_i, y_i))$ is a rung. Let $L = (\bigcup_{i=1}^m R_i) + S$. Then $(L, (x_0, v_0, y_0), (x_m, v_m, y_m))$ is a ladder.

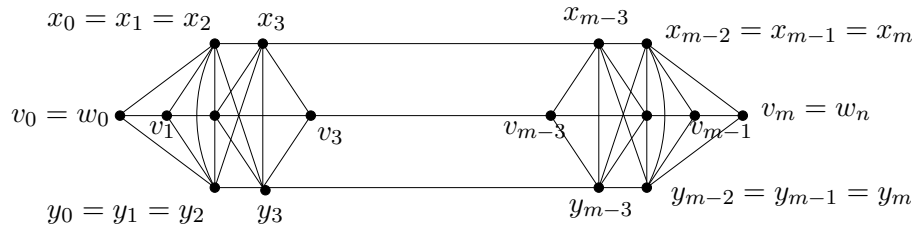


Figure 3: The ladder L .

Finally, let G be obtained from the disjoint union of J and L by identifying the vertices w_0, \dots, w_n in both J and L . Now (J, L) is a separation in G , $V(J \cap L) = \{w_0, \dots, w_n\}$, (J, w_0, \dots, w_n) is 3-planar, and $(L, (x_0, v_0, y_0), (x_m, v_m, y_m))$ is a ladder along $v_0v_1 \dots v_m$,

where $w_0 \dots w_n$ is the reduced sequence of $v_0 \dots v_m$.

By (1.3), $(G, \{x_0, y_0\}, \{x_m, y_m\}, (v_0, v_m))$ is an obstruction.

We claim that G is 7-connected. Suppose on the contrary that G is not 7-connected. Then G contains a cut set T with $|T| \leq 6$. Let X be a component of $G - T$. If $V(X) \cap \{w_0, \dots, w_n\} = \emptyset$, then $X \subseteq L - \{w_0, \dots, w_n\}$ (since $|T| \leq 6$ and by (3)), and so, $|T| \geq 7$ (by the construction of L), a contradiction. So $V(X) \cap \{w_0, \dots, w_n\} \neq \emptyset$.

If $(V(X) \cap V(L)) - \{w_0, \dots, w_n\} \neq \emptyset$, then $|T \cap V(L)| \geq 7$ (by the construction of L), a contradiction.

Hence, $V(X) \cap V(L) \subseteq \{w_0, \dots, w_n\}$, and so, X is a component of $J - (T \cap V(J))$ and the neighbors of $V(X) \cap \{w_0, \dots, w_n\}$ in L must be contained in T . By (3), $|T \cap V(J)| \geq 3$. If $V(X) \cap \{w_2, \dots, w_{n-2}\} \neq \emptyset$, then $|(T \cap L) - \{w_0, \dots, w_n\}| \geq 4$, and so, $|T \cap V(J)| \leq 2$, a contradiction. So assume that $V(X) \cap V(L) \subseteq \{w_0, w_1, w_{n-1}, w_n\}$. Furthermore, $V(X) \cap V(L) \subseteq \{w_0, w_1\}$ or $V(X) \cap V(L) \subseteq \{w_{n-1}, w_n\}$; as otherwise, $|(T \cap V(L)) - \{w_0, \dots, w_n\}| \geq 4$, and so, $|T \cap V(J)| \leq 2$, a contradiction. So there is some $x \in V(b_2Cb_1)$ such that every path from b_2Cx to xCb_1 in $J - b_1b_2$ intersects $T \cap V(J)$. Since $|(T \cap V(L)) - \{w_0, \dots, w_n\}| \geq 2$, and so, $|T \cap V(J)| \leq 4$, contradicting (4). \square

Seymour pointed out that (1.3) can be used to solve the following problem: Let G be a graph and $\{b, b', x, y\} \subseteq V(G)$; determine if G contains a path from b to b' and through x and y in that order.

To see this, we consider two cases. If $xy \in E(G)$, then this problem reduces to the problem for finding a $(\{b, x\}, \{b', y\})$ -linkage in G , and so, can be solved using (2.1). So assume that $xy \notin E(G)$. Let G' be the graph obtained from G by replacing x with vertices a', c' and joining a', c' to all neighbors of x , and replacing y with vertices a, c and joining a, c to all neighbors of y . It is easy to verify that G contains a path from b to b' through x and y in that order iff G' contains three disjoint paths from $\{a, b, c\}$ to $\{a', b', c'\}$ with no path from b to b' .

Another problem mentioned to me by Seymour is to characterize the connected graphs containing cycles through four given vertices in a prescribed order. Using a similar technique as above, this problem can be formulated as follows. Let G be a graph, $\{a_1, a_2, b_1, b_2\} \subseteq V(G)$, and $\{c_1, c_2, d_1, d_2\} \subseteq V(G)$. Characterize those G such that, for any four disjoint paths from $\{a_1, a_2, b_1, b_2\}$ to $\{c_1, c_2, d_1, d_2\}$, the two paths with an end in $\{a_1, a_2\}$ have their other ends both in $\{c_1, c_2\}$ or both in $\{d_1, d_2\}$. The planar case is solved in ([4], Theorem 2.1).

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